

Trusted Routing for VANET

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- Introduction:
 - Trust in VANET
 - Challenge
 - Our approach
- Assumptions
- Proposed framework
- Applying the framework to OLSR
- Evaluation of the framework
 - Resilience to attacks
 - Computational and bandwidth overhead
- Conclusion and further work



- Trust in VANET
 - Cooperative nature: Vulnerable
 - Lack of trust standard in VANET routing protocol
- Challenges of trust establishment for VANET
 - Highly dynamic
 - Distributed
 - Resource constraints
- Trusted routing framework
 - Authentication of messages, nodes and routes
 - Limited assistance with off-line Certificate Authority (CA)

- Three-module framework:
 - Message authentication
 - Node-to-node authentication
 - Cumulative routability verification
- Prerequisites:
 - All nodes are loosely synchronized (NTP/GPS)
 - Each node has generated a key pair, K_i^+ / K_i^-
 - Off-line CA distributes following components to each node
 - Public key of CA, K_{ca}^+
 - Certificate that binds its network ID (e.g. IP address) and its public key: $Cert_i = [ID_i, K_i^+, T_v, T_e]_{K_{ca}^-}$

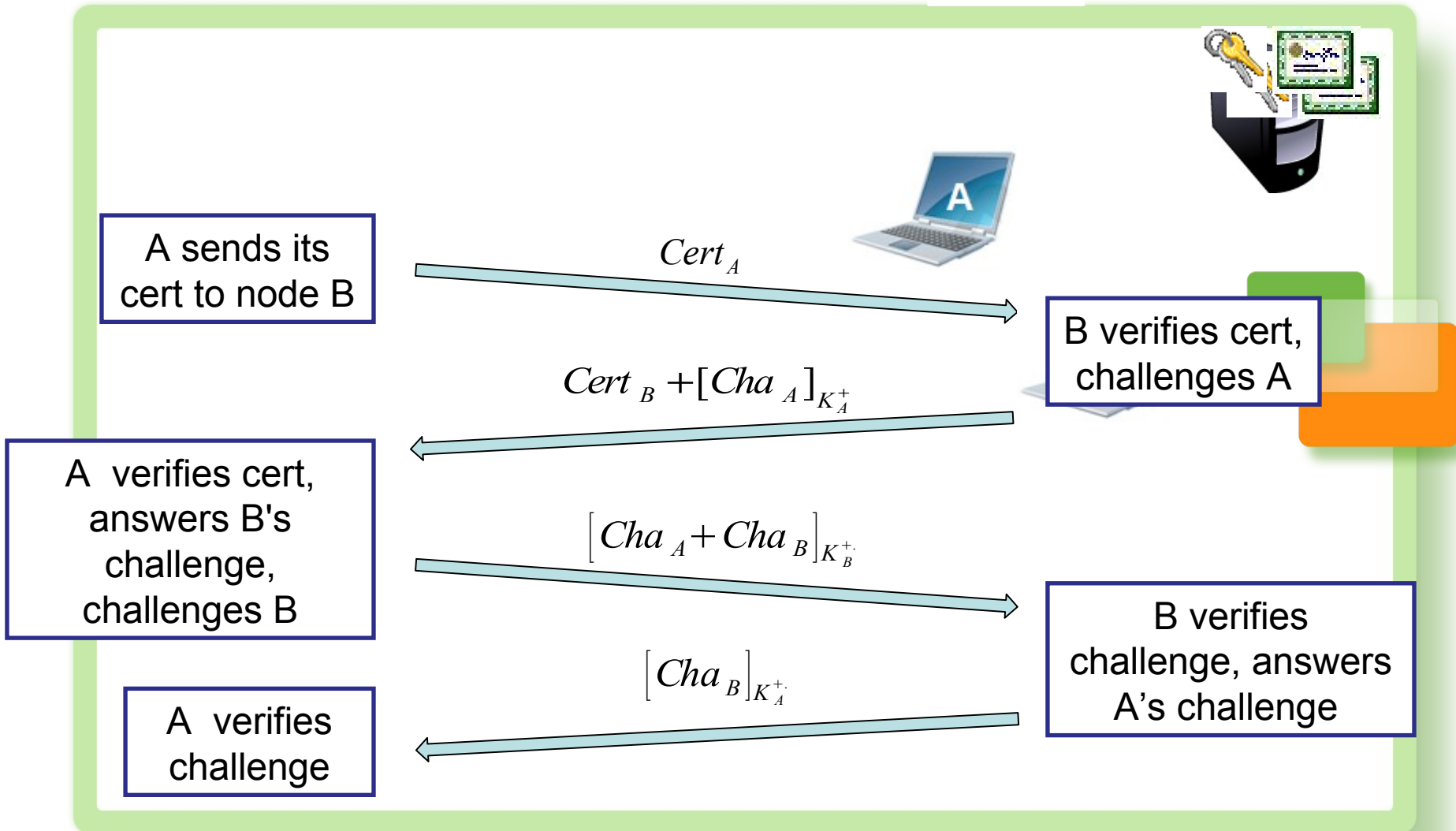
- Purpose:
 - To protect routing control messages
- Originator digitally signs every message
 - Message integrity
 - Message authentication
 - Non-repudiation
- Do not include variable fields
 - Hop-count
 - Time-to-live



- Purpose:
 - Defines a way to verify nodes in minimum iterations
- Authentication between two nodes
- Exchange certificate & public keys
- Challenge peer to confirm identity
 - i.e. possession of corresponding private key
- Exchange secret keys for quick re-authentication



Node-to-node authentication (cont.)



- Purpose:
 - Verify hop-by-hop connectivity along path
- A node must provide a piece of evidence to prove the connectivity
- Evidence: Routability Certificate (RC)
 - Signature from neighboring node regarding to the link
 - Exchange RCs after node-to-node authentication
 - Originator uses RCs to prove connectivity
 - Verify a route cumulatively

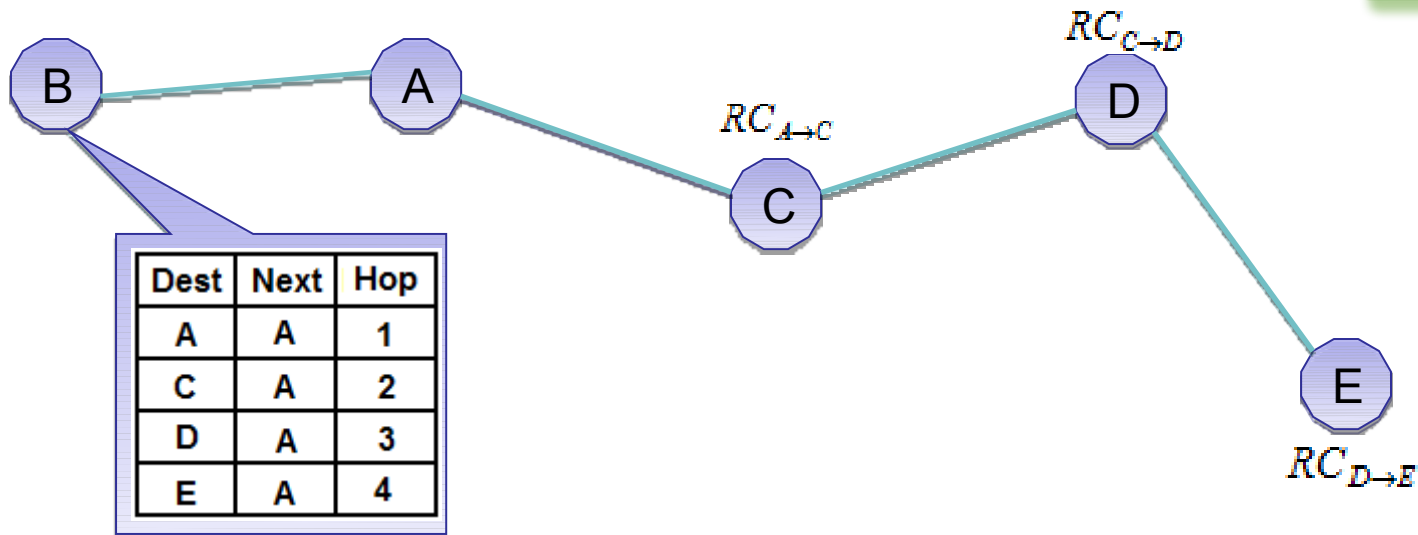


Cumulative Routability Verification (cont.)

$$RC = [ID_A, ID_B, T_v, T_e, Sign_B]$$

Where:

$$Sign_B = H([ID_A, T_v, T_e])_{K_B^-}$$

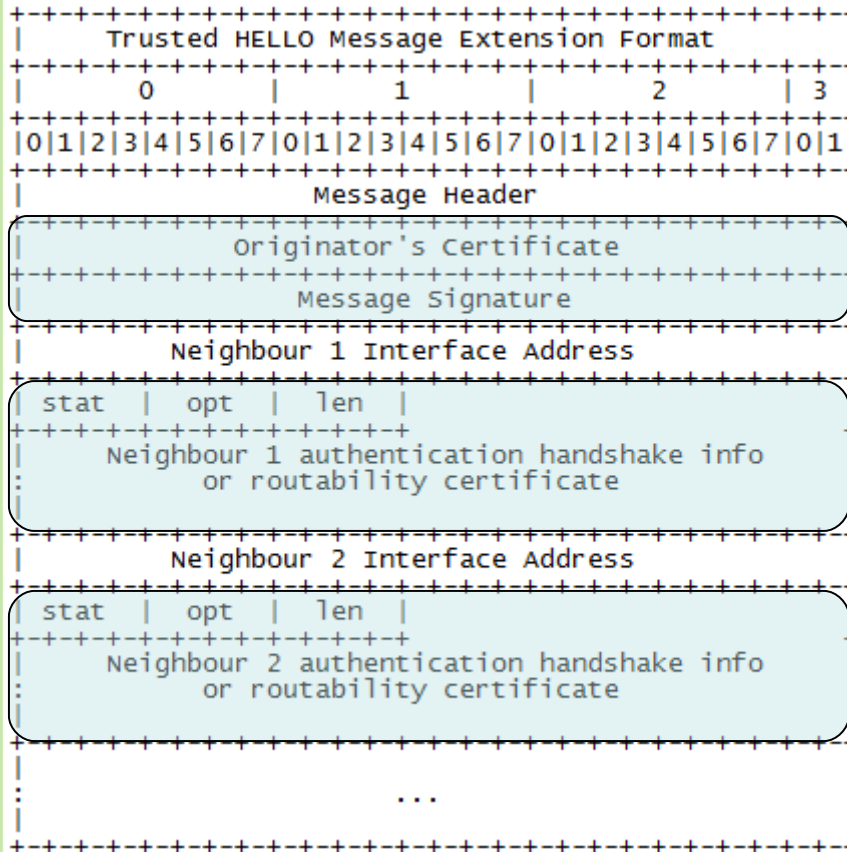


- Optimised Link State Routing (OLSR) protocol
 - Table driven, proactive
 - Use Multipoint Relays (MPR) to reduce control messages
 - Link status is disseminated to the entire network
 - HELLO message
 - Local control message
 - Link sensing
 - Neighbor discovery
 - MPR selector set discovery
 - Topology Control (TC) message
 - Global control message
 - Link state announcement

- HELLO message extension
 - Message signature
 - Authentication info is embedded in standard HELLO messages
 - Concurrent handshakes among multiple neighbors
- Operations:
 - Distribute certificates
 - Node-to-node authentication handshake
 - Exchange RCs



- Trusted HELLO message format



- Certificates
- Message Signature
- Handshake info
- RCs



Trusted Extension for OLSR (cont.)

- TC message extensions
 - Message signature
 - Carry RC and certificate after each neighbor address
 - Similar format to trusted HELLO message
- Operation:
 - Verify RC before add links to routing table
 - Confirm connection to each node hop-by-hop



Resilience to Attacks

Attack/ misbehavior	Description	Countermeasure
Illegal Access	Device without permission/certificate	Message signature Node-to-node authentication
Impersonation	Identity spoofing: MAC or IP	Node-to-node authentication
Message modification	Rushing ANSN attack	Message signature
Link spoofing	Spoofing destination that couldn't reach	Routability verification
MPR selector isolation	A node isolate its MPR selector by not include it in the message	Routability certificate sign by all neighbors

- Public key scheme dependent
- Bandwidth overhead

Size of RC:

$$L_{rc} = 2 \times L_{ip} + L_{time} + L_{sig}$$

Size of Certificate:

$$L_{cert} = L_{ip} + L_{pub} + L_{time} + L_{sig}$$

- Benchmark for some cryptographic algorithm

Operation	Milliseconds/Operation
RSA 1024 Encryption / Decryption	0.08 / 1.46
DSA 1024 Signature / Verification	0.45 / 0.52
RSA 2048 Encryption / Decryption	0.16 / 6.08
RSA 2048 Signature / Verification	6.05 / 0.16
ECIES 233 Encryption / Decryption	21.17 / 12.15
ECDSA 233 Signature / Verification	10.62 / 12.80
MD5	0.0045 (per 1KB data)
SHA-1	0.0065 (per 1KB data)

Crypto++ 5.6.0, Intel Core2 Duo 1.83 GHz

- Proposed trust establishment framework
 - Message authentication
 - Node-to-node authentication
 - Routability Verification
- Future work
 - Find a signature scheme to reduce overhead
 - Apply framework to other protocol, e.g. AODV





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The logo for NICTA (National ICT Research Australia) is centered in a white rounded square. It features a stylized circular graphic composed of a white circle on the left and a dark blue circle on the right, partially overlapping. Below the graphic, the word "NICTA" is written in a bold, green, sans-serif font.

From imagination to **impact**